

Extremal jumps of circuit complexity of unitary evolutions generated by random Hamiltonians

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Unitary complexity

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- ▶ State complexity: fix an initial state $|0\rangle$ and let $C_\epsilon(|\psi\rangle)$ be the smallest k such that for some $V_1, \dots, V_k \in \mathcal{G}$ we have:

$$d_{tr}(|\psi\rangle, V_1 \cdot V_2 \cdot \dots \cdot V_k |0\rangle) < \epsilon$$

(trace distance)

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- ▶ AdS-CFT correspondence: **Brown-Susskind conjecture** – complexity of a chaotic Hamiltonian evolution (or some toy model thereof) typically grow linearly with time t for exponentially long time.

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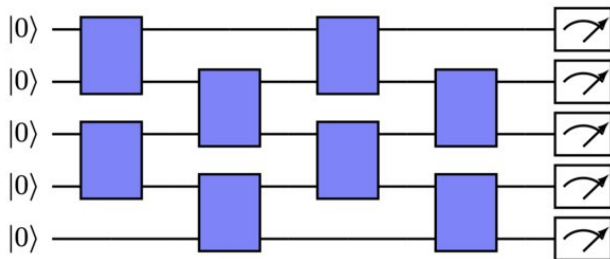
Note: lower bounding complexity for *specific* unitaries seems hopelessly difficult! Classical circuit lower bounds – $P \neq NP$ type of questions.

Random Quantum Circuits

So, time to introduce some randomness...

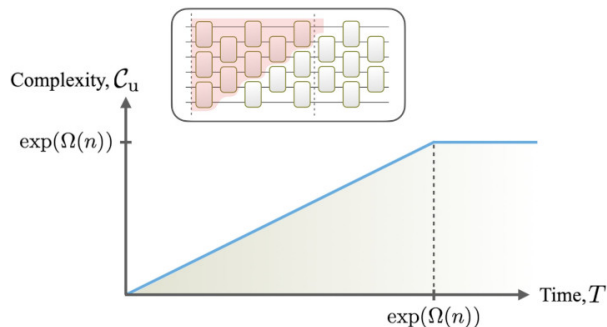
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So, time to introduce some randomness...



- ▶ choose local gates $U_{i,i+1}$ to be Haar random unitaries from 2-qubit unitaries $U(4)$, run for depth t .
- ▶ approaches Haar measure on $U(d)$ as $t \rightarrow \infty$; k -design properties known (Brandao, Harrow, Horodecki 2016).

Random Quantum Circuits



- ▶ **Brown-Susskind conjecture:** complexity of a random quantum circuit typically grows linearly with depth t until exponential value (no shortcuts). Proved in some settings (exact complexity) by Haferkamp, Faist, Kothakonda, Eisert, Younger Halpern (2022) and Haferkamp (2023). *Still open!*

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If complexity $> k$, then U lies outside the union of ε -balls around $\{G_1, G_2, \dots\} \in \mathcal{G}^k$ (words of length at most k)...

$$B(G_3, \varepsilon) \quad B(G_5, \varepsilon)$$

$$B(G_1, \varepsilon) \quad B(I, \varepsilon) \quad B(G_4, \varepsilon)$$

$$B(G_2, \varepsilon)$$

Unitary complexity

A single ball has volume $\sim \varepsilon^{d^2}$, so total volume $\leq |\mathcal{G}|^k \varepsilon^{d^2}$ – very small if $k \lesssim \frac{1}{\log |\mathcal{G}|} \cdot d^2 \log \left(\frac{1}{\varepsilon} \right)$

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Solovay-Kitaev theorem: every U has complexity at most $c(\mathcal{G}) \cdot d^2 \log^\gamma \left(\frac{1}{\varepsilon}\right)$ for some $\gamma < 4$. So *Haar random unitaries have typically maximal complexity.*

Our results

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Would like to prove something for **local** Hamiltonian with random couplings:

$$H = \sum_i g_i H_i$$

where H_i - local, $\|H_i\| = 1$, constants g_i - random couplings.

Difficult :(

Our results

Instead, consider a "mean-field" random matrix GUE model (prominent e.g. in quantum chaos):

- ▶ The random Hamiltonian H is taken to be the $d \times d$ GUE matrix (Gaussian Unitary Ensemble, standard in random matrix theory, strongly interacting) – independent Gaussian entries:

$$H_{ii} \sim \text{real Gaussian of variance } 1$$

$$H_{ij} \sim \text{complex Gaussian of variance } 1$$

$$H_{ij} = \overline{H_{ji}}$$

- ▶ caveat: H is nonlocal

Our results

Theorem (Complexity jump for GUE evolutions)

Let $t_{\text{esc}} = C\varepsilon$, $t_{\text{jump}} = C'\varepsilon$ ($C' > C$). With high probability over $H \sim \text{GUE}(d)$

- (a) For all times $t \in [0, t_{\text{esc}}]$ complexity of evolution U_t is trivial: $C_\varepsilon(U_t) = 0$.
- (b) For any fixed time $t > t_{\text{jump}}$ complexity satisfies $C_\varepsilon(U_t) \geq \frac{C''\varepsilon^2}{\log |G|} d^2$ (almost maximal wrt. d)

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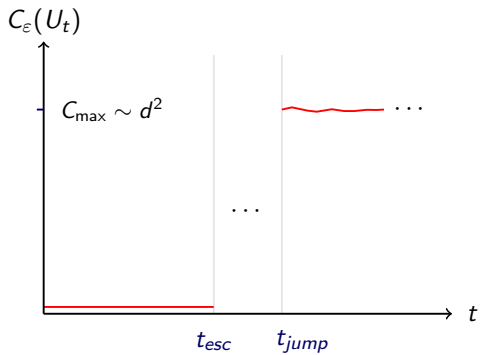
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Much more powerful than non-rigorous heuristics based on frame potentials and moment computations (Cotler, Hunter-Jones, Liu, Yoshida 2017).

Our results



Very high level proof ideas

- Idea 1: GUE ensemble is unitarily invariant: the eigenbasis of H is Haar random! Use it to exclude all balls around $G \neq I$ (technically: concentration of measure on the unitary group)

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- Idea 2: use concentration of measure for the **spectrum** of GUE matrix to control how spread out the distribution is (not too concentrated on the ball around identity).

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Thank you!